### More SQL

#### Extended Relational Algebra Outerjoins, Grouping/Aggregation

### The Extended Algebra

- $\delta$  = eliminate duplicates from bags.
- T =sort tuples.
- $\gamma$  = grouping and aggregation.

*Outerjoin* : avoids "dangling tuples" = tuples that do not join with anything.

### **Duplicate Elimination**

#### R1 := δ(R2).

#### R1 consists of one copy of each tuple that appears in R2 one or more times.

### **Example:** Duplicate Elimination

$$R = ( \begin{array}{c|c} A & B \\ 1 & 2 \\ 3 & 4 \\ 1 & 2 \end{array} )$$

<b>δ</b> (R)	=	

# Sorting

#### • R1 := $T_{L}$ (R2).

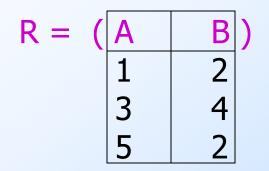
L is a list of some of the attributes of R2.

R1 is the list of tuples of R2 sorted first on the value of the first attribute on *L*, then on the second attribute of *L*, and so on.

Break ties arbitrarily.

T is the only operator whose result is neither a set nor a bag.

### **Example:** Sorting



 $T_B(R) = [(5,2), (1,2), (3,4)]$ 

### **Aggregation Operators**

 Aggregation operators are not operators of relational algebra.

- Rather, they apply to entire columns of a table and produce a single result.
- The most important examples: SUM, AVG, COUNT, MIN, and MAX.

### **Example:** Aggregation

$$R = (A B) \\
1 3 \\
3 4 \\
3 2$$

SUM(A) = 7COUNT(A) = 3MAX(B) = 4AVG(B) = 3

# **Grouping Operator**

• R1 :=  $\gamma_{L}$  (R2). L is a list of elements that are either:

- 1. Individual (*grouping* ) attributes.
- 2. AGG(*A*), where AGG is one of the aggregation operators and *A* is an attribute.
  - An arrow and a new attribute name renames the component.

# Applying $\gamma_{L}(R)$

- Group *R* according to all the grouping attributes on list *L*.
  - That is: form one group for each distinct list of values for those attributes in *R*.
- Within each group, compute AGG(A) for each aggregation on list L.
  - Result has one tuple for each group:
    - 1. The grouping attributes and
    - 2. Their group's aggregations.

# **Example:** Grouping/Aggregation

$$R = \left(\begin{array}{c|cc} A & B & C \\ 1 & 2 & 3 \\ 4 & 5 & 6 \\ 1 & 2 & 5 \end{array}\right)$$

 $Y_{A,B,AVG(C)->X}(R) = ??$ 

First, group *R* by *A* and *B*:

A	B	С
1	2	3
1	2	5
4	5	6

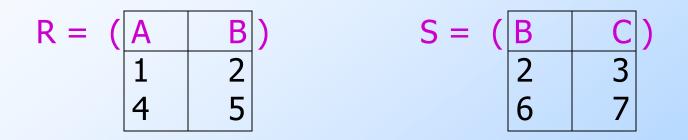
Then, average *C* within groups:

Α	В	Х
1	2	4
4	5	6

# Outerjoin

- ◆ Suppose we join *R* ⋈<sub>C</sub> *S*.
   ◆ A tuple of *R* that has no tuple of *S* with which it joins is said to be *dangling*.
  - Similarly for a tuple of S.
- Outerjoin preserves dangling tuples by padding them NULL.

### **Example:** Outerjoin



(1,2) joins with (2,3), but the other two tuples are dangling.

### Now --- Back to SQL

#### Each Operation Has a SQL Equivalent

# Outerjoins

- R OUTER JOIN S is the core of an outerjoin expression. It is modified by:
  - 1. Optional NATURAL in front of OUTER:
  - 2. Optional ON < condition > after JOIN.
  - Optional LEFT, RIGHT, or FULL before OUTER.
     Only one
    - LEFT = pad dangling tuples of R only.
    - RIGHT = pad dangling tuples of S only.
    - FULL = pad both; this choice is the default.

of these

### Aggregations

SUM, AVG, COUNT, MIN, and MAX can be applied to a column in a SELECT clause to produce that aggregation on the column.

 Also, COUNT(\*) counts the number of tuples.

### **Example:** Aggregation

◆From Sells(bar, beer, price), find the
average price of Bud:
 SELECT AVG(price)
 FROM Sells
 WHERE beer = 'Bud';

# Eliminating Duplicates in an Aggregation

Use DISTINCT inside an aggregation.
 Example: find the number of different prices charged for Bud:
 SELECT COUNT(DISTINCT price)
 FROM Sells
 WHERE beer = 'Bud';

# NULL's Ignored in Aggregation

 NULL never contributes to a sum, average, or count, and can never be the minimum or maximum of a column.

 But if there are no non-NULL values in a column, then the result of the aggregation is NULL.

Exception: COUNT of an empty set is 0.

### Example: Effect of NULL's

SELECT count(\*) FROM Sells WHERE beer = 'Bud';

The number of bars that sell Bud.

SELECT count(price) FROM Sells WHERE beer = 'Bud';

The number of bars that sell Bud at a known price.

### Grouping

 We may follow a SELECT-FROM-WHERE expression by GROUP BY and a list of attributes.

The relation that results from the SELECT-FROM-WHERE is grouped according to the values of all those attributes, and any aggregation is applied only within each group.

### **Example:** Grouping

◆From Sells(bar, beer, price), find the average price for each beer: SELECT beer, AVG(price) FROM Sells GROUP BY beer;

beer	AVG(price)
Bud	2.33

# **Example:** Grouping

From Sells(bar, beer, price) and Frequents(drinker, bar), find for each drinker the average price of Bud at the bars they frequent:

SELECT drinker, AVG(price)

FROM Frequents, Sells

**GROUP BY drinker;** 

WHERE beer = 'Bud' AND

Frequents.bar = Sells.bar

Compute all drinker-barprice triples for Bud.

Then group them by drinker.

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# Restriction on SELECT Lists With Aggregation

- If any aggregation is used, then each element of the SELECT list must be either:
  - 1. Aggregated, or
  - 2. An attribute on the GROUP BY list.

# **Illegal Query Example**

You might think you could find the bar that sells Bud the cheapest by: SELECT bar, MIN(price) FROM Sells WHERE beer = 'Bud';

But this query is illegal in SQL.

### **HAVING Clauses**

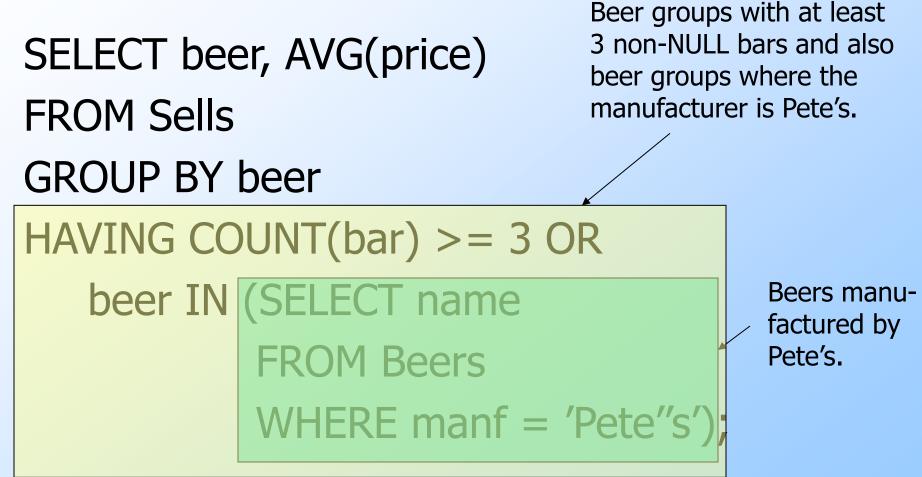
 HAVING <condition> may follow a GROUP BY clause.

 If so, the condition applies to each group, and groups not satisfying the condition are eliminated.

### **Example: HAVING**

From Sells(bar, beer, price) and Beers(name, manf), find the average price of those beers that are either served in at least three bars or are manufactured by Pete's.

### Solution



# Requirements on HAVING Conditions

- Anything goes in a subquery.
- Outside subqueries, they may refer to attributes only if they are either:
  - 1. A grouping attribute, or
  - 2. Aggregated

(same condition as for SELECT clauses with aggregation).