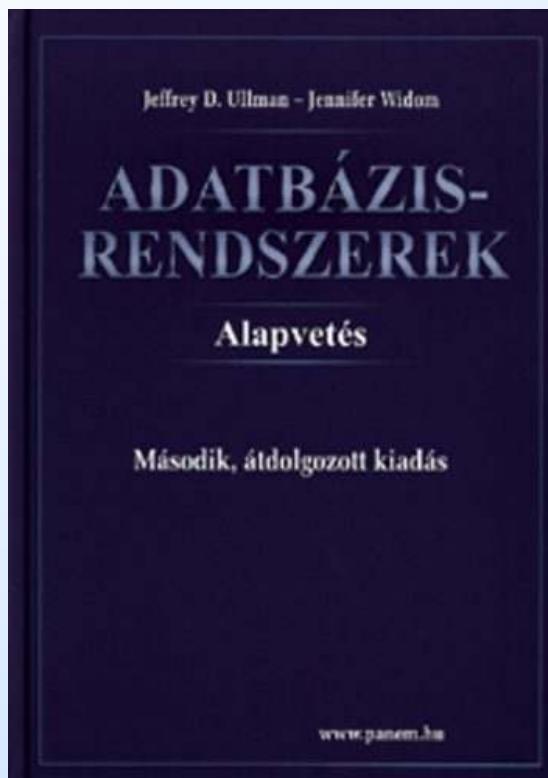


Datalog -2: Rekurzió



Ullman-Widom: Adatbázisrendszerek
Alapvetés

Második, átdolgozott kiadás, Panem,
2009

10.2. Rekurzió Datalogban és
SQL-99-ben

(Jeffrey D. Ullman, 2007)

Datalog

- ◆ Datalog program – véges sok Datalog szabály megadásából áll
- ◆ Extenzionális, intenzionális predikátumok
- ◆ Datalog program kiértékelése
- ◆ Rekurzió megadása, ha nem használunk negációt

Expressive Power of Datalog

- ◆ Without recursion, Datalog can express all and only the queries of core relational algebra.
 - ▷ The same as SQL select-from-where, without aggregation and grouping.
- ◆ But with recursion, Datalog can express more than these languages.

Recursive Example

- ◆ EDB: $\text{Par}(c,p) = p$ is a parent of c .
- ◆ Generalized cousins: people with common ancestors one or more generations back:

$\text{Sib}(x,y) \leftarrow \text{Par}(x,p) \text{ AND } \text{Par}(y,p) \text{ AND } x <> y$

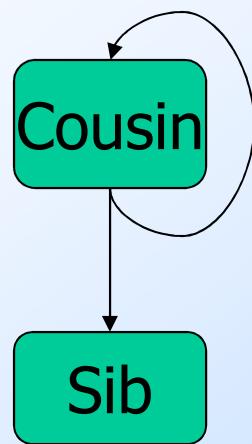
$\text{Cousin}(x,y) \leftarrow \text{Sib}(x,y)$

$\text{Cousin}(x,y) \leftarrow \text{Par}(x,xp) \text{ AND } \text{Par}(y,yp)$
AND $\text{Cousin}(xp,yp)$

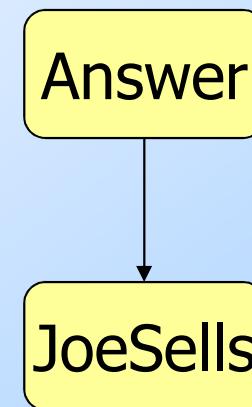
Rekurzió definíciója

- ◆ Megelőzési gráf pontjai P1, ... Pn
élek $P_i \rightarrow P_j$ -be, ha van olyan szabály,
amelynek feje P_j és törzsében szerepel P_i
- ◆ Form a *dependency graph* whose
nodes = IDB predicates.
- ◆ Arc $X \rightarrow Y$ if and only if there is a rule
with X in the head and Y in the body.
- ◆ Cycle = recursion; no cycle = no recursion

Example: Dependency Graphs



Recursive

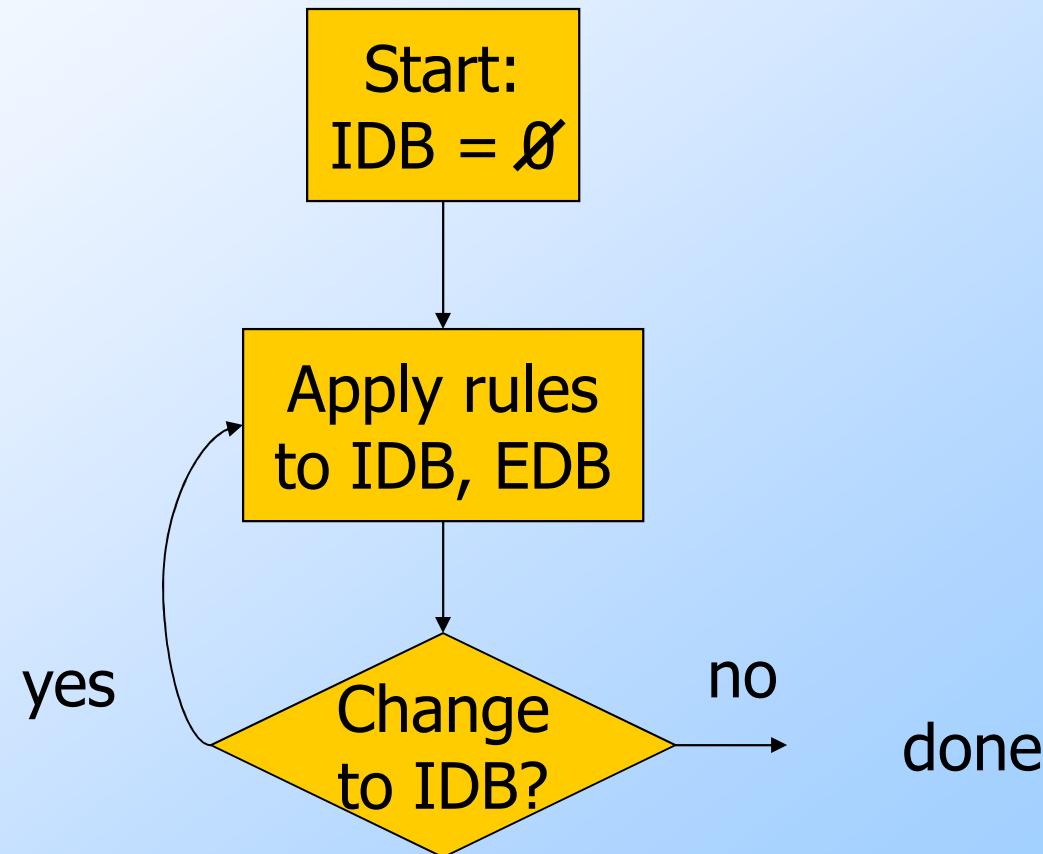


Nonrecursive

Evaluating Recursive Rules

- ◆ The following works when there is no negation:
 1. Start by assuming all IDB relations are empty.
 2. Repeatedly evaluate the rules using the EDB and the previous IDB, to get a new IDB.
 3. End when no change to IDB.

The “Naïve” Evaluation Algorithm



Seminaive Evaluation

- ◆ Since the EDB never changes, on each round we only get new IDB tuples if we use at least one IDB tuple that was obtained on the previous round.
- ◆ Saves work; lets us avoid rediscovering *most* known facts.
 - ▷ A fact could still be derived in a second way.

SQL-99 Recursion

- ◆ Datalog recursion has inspired the addition of recursion to the SQL-99 standard. (IBM DB2 implementálta)
- ◆ Tricky, because SQL allows negation grouping-and-aggregation, which interact with recursion in strange ways.
- ◆ Megszorítások a rekurzióra:
 - monoton lekérdezések
 - csak lineáris rekurzió szerepel

Form of SQL Recursive Queries

WITH

[RECURSIVE] <relnév1> AS (SELECT...)

[RECURSIVE] <relnév2> AS (SELECT...)

...

[RECURSIVE] <relnévk> AS (SELECT...)

SELECT ... FROM listán <relnévi>-k

Example: SQL Recursion – (1)

- ◆ Find Sally's cousins, using SQL like the recursive Datalog example.
- ◆ Par(child,parent) is the EDB.

WITH Sib(x,y) AS

```
SELECT p1.child, p2.child  
FROM Par p1, Par p2  
WHERE p1.parent = p2.parent AND  
      p1.child <> p2.child;
```

Like Sib(x,y) <-
Par(x,p) AND
Par(y,p) AND
 $x \neq y$

Example: SQL Recursion – (2)

```
WITH ...          Required – Cousin  
    is recursive  
RECURSIVE Cousin(x,y) AS  
    (SELECT * FROM Sib)          Reflects Cousin(x,y) <-  
                                Sib(x,y)  
    UNION  
    (SELECT p1.child, p2.child  
     FROM Par p1, Par p2, Cousin  
     WHERE p1.parent = Cousin.x AND  
           p2.parent = Cousin.y);          Reflects  
                                         Cousin(x,y) <-  
                                         Par(x,xp) AND  
                                         Par(y,yp) AND  
                                         Cousin(xp,yp)
```

Example: SQL Recursion – (3)

- ◆ With those definitions, we can add the query, which is about the virtual view Cousin(x,y):

```
SELECT y  
FROM Cousin  
WHERE x = 'Sally' ;
```

Legal SQL Recursion

- ◆ It is possible to define SQL recursions that do not have a meaning.
- ◆ The SQL standard restricts recursion so there is a meaning.
- ◆ And that meaning can be obtained by seminaïve evaluation.

Example: Meaningless Recursion

- ◆ EDB: $P(x) = \{(1)\}$.
- ◆ IDB: $Q(x) \leftarrow P(x) \text{ AND NOT } Q(x)$.
- ◆ Is (1) in $Q(x)$?
 - ▶ If so, the recursive rule says it is not.
 - ▶ If not, the recursive rule says it is.

Plan to Explain Legal SQL Recursion

1. Define “monotone” recursions.
2. Define a “stratum graph” to represent the connections among subqueries.
3. Define proper SQL recursions in terms of the stratum graph.

Monotonicity

- ◆ If relation P is a function of relation Q (and perhaps other relations), we say P is *monotone* in Q if inserting tuples into Q cannot cause any tuple to be deleted from P .
- ◆ Examples:
 - ▶ $P = Q \cup R$.
 - ▶ $P = \sigma_{a=10}(Q)$.

Rekurzió a gyakorlatban

- ◆ SQL-99 szabvány csak az ún. „monoton” rekurziót támogatja (viszont Datalogban megengedett a negáció és rekurzió együtt, de igen bonyolulttá válik a lekérdezés értelmezése, mint például a rétegzés),
- ◆ SQL-99 szabvány lineáris rekurziót enged meg, vagyis az értékkadás kifejezésben egyetlen rekurzív relációt használhatunk.

Tankönyv példája

- ◆ Jaratok(legitarsasag, honnan, hova, koltseg, indulas, erkezes) táblában repülőjáratok adatait tároljuk.
- ◆ Mely (x,y) párokra lehet eljutni x városból y városba?
- ◆ Eljut(x, y) <- Jaratok(_, x, y, _, _, _)
Eljut(x, y) <- Eljut(x, z) AND Jaratok(_, z, y, _, _, _)
- ◆ WITH RECURSIVE Eljut AS
(SELECT honnan, hova FROM Jaratok
UNION
(SELECT Eljut.honnan, Jaratok.hova
FROM Eljut, Jaratok
WHERE Eljut.hova = Jaratok.honnan)
SELECT * FROM eljut;
- ◆ Lásd még a gyakorlaton is PL/SQL programmal kifejezve! ²⁰