# Single Table, Index (2.1)

SELECT \* FROM emp WHERE empno=174;

Unique emp(empno)

>.SELECT STATEMENT

- >...TABLE ACCESS by rowid emp
- >....INDEX unique scan i emp pk

- Index Unique Scan
  - Traverses the node blocks to locate correct leaf block
  - Searches value in leaf block (if not found => done)
  - Returns rowid to parent row-source
    - · Parent: accesses the file+block and returns the row

# Index Unique Scan (2.1)



# Single Table, Index (2.2)



- (Non-unique) Index Range Scan
  - Traverses the node blocks to locate most left leaf block
  - Searches 1<sup>st</sup> occurrence of value in leaf block
  - Returns rowid to parent row-source
    - Parent: accesses the file+block and returns the row
  - Continues on to next occurrence of value in leaf block
    - Until no more occurences

#### Index Range Scan (2.2)



# Single Table, Index (2.3)

#### SELECT \* FROM emp WHERE empno>100;

Unique emp(empno)

- >.SELECT STATEMENT
- >...TABLE ACCESS by rowid emp
- >....INDEX range scan i\_emp\_pk

- Unique Index Range Scan
  - Traverses the node blocks to locate most left leaf block with start value
  - Searches 1<sup>st</sup> occurrence of value-range in leaf block
  - Returns rowid to parent row-source
    - Parent: accesses the file+block and returns the row
    - Continues on to next valid occurrence in leaf block
      - Until no more occurences / no longer in value-range



Multiple levels of Btrees, by column order

# Single Table, Index (2.4)

#### **SELECT** \*

FROM emp WHERE job='manager' AND hiredate='01-01-2001';

```
Emp(job,hiredate)
```

- >.SELECT STATEMENT
- >...TABLE ACCESS by rowid emp
- >....INDEX range scan i\_emp\_j\_h

- Full Concatenated Index
  - Use job-value to navigate to sub-Btree
  - Then search all applicable hiredates

## Single Table, Index (2.5)

| CEI |   | * |
|-----|---|---|
| SEI | ノ |   |

FROM emp WHERE job='manager';

Emp(job,hiredate)

>.SELECT STATEMENT

- >...TABLE ACCESS by rowid emp
- >....INDEX range scan i\_emp\_j\_h

- (Leading) Prefix of Concatenated Index
  - Scans full sub-Btree inside larger Btree

### Index Range Scan (2.5)





# Single Table, Index (2.6)



- >.SELECT STATEMENT
- >...TABLE ACCESS by rowid emp
- >....INDEX range scan i\_emp\_j\_h

- Index Skip Scan (prior versions did FTS)
  - "To use indexes where they've never been used before"
  - Predicate on leading column(s) no longer needed
  - Views Btree as collection of smaller sub-Btrees
  - Works best with low-cardinality leading column(s)



# Single Table, Index (2.7)

>.SELECT STATEMENT

>...TABLE ACCESS by rowid emp

>....INDEX range scan i emp job

#### **SELECT** \*

FROM emp WHERE empno>100 AND job='manager';

```
Unique Emp(empno)
Emp(job)
```

#### • Multiple Indexes

- Rule: uses heuristic decision list to choose which one
  - Avaliable indexes are 'ranked'
- Cost: computes most selective one (ie. least costing)
  - Uses statistics

# **RBO** Heuristics

- Ranking multiple available indexes
  - 1. Equality on single column unique index
  - 2. Equality on concatenated unique index
  - 3. Equality on concatenated index
  - 4. Equality on single column index
  - 5. Bounded range search in index
    - Like, Between, Leading-part, ...
  - 6. Unbounded range search in index
    - Greater, Smaller (on leading part)

Normally you hint which one to use

## **CBO** Cost Computation

- Statistics at various levels
  - Table:
    - Num\_rows, Blocks, Empty\_blocks, Avg\_space
  - Column:
    - Num\_values, Low\_value, High\_value, Num\_nulls
  - Index:
    - Distinct\_keys, Blevel, Avg\_leaf\_blocks\_per\_key, Avg\_data\_blocks\_per\_key, Leaf\_blocks
  - Used to compute selectivity of each index
    - Selectivity = percentage of rows returned
      - Number of I/O's plays big role
    - FTS is also considered at this time!